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PATENT APPLICATION

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

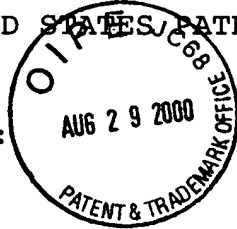
In re Application of:

PAWEL REJ

Application No.: 09/599,950

Filed: June 23, 2000

For: SPLIT TREE DATA STRUCTURE) August 29, 2000



: Examiner: Unassigned

: Group Art Unit: Unassigned

The Commissioner for Patents  
Washington, D.C. 20231

CLAIM TO PRIORITY

Applicant hereby claims priority under the  
International Convention and all rights to which he is entitled  
under 35 U.S.C. § 119 based upon the following Australian  
Priority Application:

42635/00

Australia

24 June 1999.

A certified copy of the priority document is  
enclosed.

Applicant's undersigned attorney may be reached in  
our Washington, D.C. office by telephone at (202) 530-1010. All

correspondence should continue to be directed to our address  
given below.

Respectfully submitted,

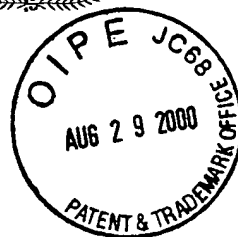


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Patent Office  
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I, LEANNE MYNOTT, TEAM LEADER EXAMINATION SUPPORT AND SALES hereby certify that annexed is a true copy of the Provisional specification in connection with Application No. PQ 1175 for a patent by CANON KABUSHIKI KAISHA filed on 24 June 1999.

I further certify that pursuant to the provisions of Section 38(1) of the Patents Act 1990 a complete specification was filed on 23 June 2000 and it is an associated application to Provisional Application No. PQ 1175 and has been allocated No. 42635/00.

## CERTIFIED COPY OF PRIORITY DOCUMENT

WITNESS my hand this  
Fifth day of July 2000

LEANNE MYNOTT  
TEAM LEADER EXAMINATION  
SUPPORT AND SALES

App/n. No.: 09/599,950

Filed: June 23, 2000

Inv: PAWEL REJ

Title: Split Tree DATA structure



**ORIGINAL**

**AUSTRALIA**

**Patents Act 1990**

**PROVISIONAL SPECIFICATION FOR THE INVENTION ENTITLED:**

Split Tree Data Structure

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Name and Address  
of Applicant:

Canon Kabushiki Kaisha, incorporated in Japan, of 30-2,  
Shimomaruko 3-chome, Ohta-ku, Tokyo, 146, JAPAN

Name of Inventor: Pawel Rej

This invention is best described in the following statement:

## **SPLIT TREE DATA STRUCTURE**

### **Technical Field of the Invention**

The present invention relates generally to document formatting and, in particular, to document formatting using input trees. More specifically, the present invention relates to a method and apparatus for creating a split tree by assigning marks to an input tree and for generating tree fragments from the split tree. The invention also relates to a computer program product including a computer readable medium having recorded thereon a computer program for creating the split tree by assigning marks to the input tree and for generating tree fragments from the split tree.

### **Background Art**

High-level languages for document formatting exist in many forms, some of which accept as input a description of a document in the form of an input tree. An example of such an input tree is presented in Fig. 1. The input tree comprises two main types of elements, namely data elements and instruction elements, those elements in Fig. 1 being indicated in Fig. 2. The data elements contain data to be formatted. Each data element can include text or images. Instruction elements contain information used by the formatting algorithm to arrange the data element on the pages when producing an output. As seen in Fig. 3, the data elements form leaf nodes in the tree whereas the instruction elements form nodes. Further, nodes A and B are child nodes of node Cat(2). The left node, Node A in the example, is defined as the first child, whereas the right node, Node B, is defined as the second child.

Instruction elements can be thought of as operators that define the output positions of data elements. There are a number of instruction elements or operators available. Certain of the instruction elements, for example those defining text size, text spacing and the like, typically only have one child.

The other relevant instruction in Figs. 1 to 3 is “Cat( )”, which performs vertical concatenation. Each concatenation instruction places its left parameter (the first child) above its right parameter (second child). For example, Cat(2) would place the content of data element B below the content of data element A.

5           As an input, data elements do not have an actual position. The actual position of the data elements is determined by a document formatter by evaluating the input tree. In the example of Fig. 1, the formatter, following the directions given by the instruction elements, determines that the result of the Cat(2) operator is placed above the result of the Cat(3) operator. The Cat(2) operator in turn places the content of data element A above  
10 data element B, and the Cat(3) operator places the content of data element C above data element D. The result is that data element A, B, C and D are stacked in that order.

          However, when the data elements are placed on a page in order to produce an output, it is sometimes found that the result is too big to fit on one page. Therefore, following the order that data elements are placed in, should the next data element not fit  
15 into the space left on the present page, the remainder of the data elements are promoted to the next page. In the example presented in Fig. 1, only data elements A, B and C are located on page 1 and the data element D is located on page 2. The result of the evaluation of the input tree is shown in Fig. 4. The formatter also positions the data elements on the pages.

20           To better explain the above, a galley and a galley target need to be defined. A galley is a part of the input tree, whereas a galley target is a placeholder for a galley. The galley is therefore not evaluated at a position in the input tree but rather at the point where it is placed into the galley target. The galley target is a stretchable area on the output area with size 0 when no galley is targeted into it. When a galley is targeted into a galley  
25 target, the galley target stretches to accommodate for the incoming galley. The galley is

unable to stretch infinitely though. There are restrictions on the maximum size of each galley target. Should it happens that the galley target reaches its maximum size before the galley is exhausted, a new galley target is found for the remainder of the content of the galley and the process of accommodating the remainder of the galley into its the new  
5 target is continued.

Placing a galley into a series of targets involves splitting the input tree into fragments, each fragment filling its target. Fig. 5 shows another example of a galley (input tree). It is assumed that the galley in the example of Fig. 5 is targeted into a sequence of targets. Only data elements A and B fit into the first galley target in the  
10 sequence by taking all the available space defined by the restrictions on the first galley target. The rest of the galley is promoted to the following galley target in the sequence.

Figs 6A and 6B show the result after a portion of the galley of Fig.5 is placed into the sequence of galley targets. The original galley tree shown in Fig. 5 is split into the two tree fragments. The first fragment, shown in Fig. 6A, is placed into the first  
15 galley target. The remainder of the galley, shown in Fig. 6B, is waiting to be placed into the following galley targets. Thus, the original galley was split into two separate smaller trees.

The arrangement described above is used in batch text formatters which produce an output from an input tree. However, this arrangement suffers from the disadvantage  
20 that once the galley is split into smaller trees, it is hard to relate the fragmented trees back to their original galley tree. This is for example necessary whenever the order or contents of the data objects are changed, thus preventing interactive formatting.

### **Disclosure of the Invention**

It is an object of the present invention to substantially overcome, or at least  
25 ameliorate, one or more disadvantages of existing arrangements.

According to a first aspect of the invention, there is provided a method of creating a split tree for representing both an input tree and at least one tree fragment, wherein said input tree comprises a plurality of nodes, said method comprising the steps of:

5       determining which of said plurality of nodes fit into a galley target; and  
       marking said nodes that fit into said galley target with a mark to create said tree  
       fragment for each galley target, said marked input tree being said split tree.

According to a second aspect of the invention, there is provided a method of splitting a split tree into at least one tree fragments, wherein each one of said tree fragments is associated with a mark and said split tree comprises at least one node marked  
10 said mark, said method comprising the steps of:

identifying said at least one node with said mark; and

creating said at least one tree fragment from said nodes marked with each of said marks.

15           According to another aspect of the invention, there is provided an apparatus for  
implementing any one of the aforementioned methods.

According to another aspect of the invention there is provided a computer program product including a computer readable medium having recorded thereon a computer program for implementing any one of the methods described above.

20 **Brief Description of the Drawings**

A number of preferred embodiments of the present invention will now be described with reference to the drawings, in which:

Fig. 1 is an example of an input tree;

Fig. 2 illustrates types of elements in the input tree;

Fig. 3 illustrates nodes, leaf nodes and edges in the input tree;



Fig. 4 shows the result of evaluation of the input tree;

Fig. 5 shows an original galley before splitting;

Figs. 6A and 6B show the result of splitting the original galley of Fig. 5 into smaller fragments;

5 Fig. 7 shows a split tree;

Fig. 8 shows the split tree of Fig. 7 after changes;

Fig. 9 is a flow diagram of a method of generating a split tree;

Fig. 10 is a flow diagram of a method of creating tree fragments from a split tree;

10 Fig. 11 is a flow diagram of a method of finding a next node;

Fig. 12 is a flow diagram of a method of building a tree fragment;

Fig. 13 is a schematic block diagram of a general purpose computer upon which the preferred embodiment of the present invention can be practiced.

### **Detailed Description including Best Mode**

15 In the preferred embodiment of the present invention, the actual splitting of the galley tree, as discussed above, does not take place. This allows the structure of the galley to remain intact. The advantage of preserving the original structure of the galley tree is that it allows interactive formatting. For instance, it makes it easier to re-send the galley into its targets after the galley or its target have been modified. The re-sending of  
20 the galley usually involves moving the splitting boundary between the tree fragments.

The preferred embodiment discloses that a galley is represented by a Split Tree data structure. The Split tree data structure allows a split tree to represent both the galley tree and the fragments of the galley tree that are the result of splitting the galley. The advantage is that instead of actually splitting the galley tree, marks are added to the tree.

The marks allow an algorithm operating on the split tree to interpret the split tree in the same way as tree fragments would be interpreted.

The split tree corresponding to the two galley fragments in the above example of Fig. 5 is presented in Fig. 7. The fragment of the galley tree that is placed into the first galley target is marked with markers '1'. The fragment that did not fit into the first galley target is marked with marker '0'. The marker '0' is a special marker that marks nodes that are still waiting to be placed into galley targets.

The split tree data structure allows each tree fragment in the split tree to be marked with a unique mark. In the preferred embodiment, integers are used as marks.

Fig. 8 shows another example of a split tree. The galley of Fig. 7 is placed into a sequence of galley targets. For this example however, the three data elements A, B as well as C can fit into the first galley target. This change can be a result of a change in the available space for the first galley target, data elements A, B and/or C.

Each node in the split tree can have two, one or no children. The nodes that have no children are also called leaf nodes. The criteria for interpreting the split tree are as follows:

- If a node is marked with a marker 'n', the whole subtree attached to that node belongs to the 'n' fragment.
- If both the left and right subtrees of a node contain nodes marked with 'n', the node belongs to the 'n' fragment. This criterion has further implications. If a node does not contain nodes marked with the same marker in its left and right subtrees, then such a node does not belong to any fragment and is referred to as an alienated node. Any marked node, therefore a node that is not an alienated node, belongs to one tree fragment only.

- If a node has only one child and the child subtree has a node marked with 'n', then the node belongs to the 'n' fragment.

A method 30 of adding split marks to a split tree is presented in Fig. 9. The purpose of the method 30 is to place the galley tree into its galley targets. More precisely,  
5 it generates marks on the galley tree, the marks defining which parts of the galley tree are placed into which galley target.

In step 1, the method 30, called GST, accepts three input parameters:

- node - The subtree defined by 'node' that is currently being pushed into a galley target;
- 10 • available\_space - The space available in the galley target for 'node'; and
- mark - The mark used to mark all nodes that are pushed into a galley target.

In step 3 it is determined whether 'node' is already marked. If 'node' is marked, it means that it has already been placed into some other target and the subtree defined by 'node' is therefore ignored. The method 30 is then terminated in step 4 after returning  
15 'available\_space'. In this case available\_space has not changed.

For unmarked nodes, the method 30 proceeds to step 6. Step 6 determines whether the subtree defined by 'node' can fit into 'available\_space'. If that subtree can fit into 'available\_space', 'node', together with all the other nodes in that subtree, are marked with 'mark' in step 8. As 'node' uses up some of 'available\_space',  
20 'available\_space' is decreased by the size of the subtree defined by 'node' in step 9. The value of 'available\_space' is returned.

However, if in step 6 the subtree defined by 'node' cannot fit into 'available\_space', then the method 30 continues by trying to split the subtree into smaller subtrees. In order to do so, in step 13 which follows, the method 30 verifies whether  
25 'node' is a leaf element (thus have no children nodes) or has been defined as not splitable.

In these cases, 'node' cannot be placed into the target, nor can it be split, enabling the smaller subtree of 'node' to fit into the target. In step 14, the data element represented by 'node' is discarded and 'node' is marked with a special 'discarded' mark.

If 'node' is determined to be splitable in step 13, step 17 which follows  
5 determines 'node' has only one child, then step 18 recursively calls the method 30 with the child node as 'node'.

Since leave nodes were filtered out by step 13 and nodes with only one child by step 17, 'node' must have two children. When this occurs step 19 recursively calls the method 30 with node's first child of as 'node', and then, if there is space left in the target,  
10 which is determined in step 24, the method 30 is called again in step 25 with node's second child as 'node'. The method 30 terminates by returning the available space in the galley target.

The method 30 can be repeated with the next galley target, until all the data elements are placed.

15 The result is that galleys in the input tree are now converted into split trees by marking the nodes. Fig. 10 presents a method 40 for generating tree fragments from a split tree. The method 40 calls procedures 50 and 70, presented in Figs 11 and 12 respectively. These procedures 50 and 70 are discussed first.

Procedure 50 is a recursive procedure that finds a next node belonging to a tree  
20 fragment in a split tree. This tree fragment is one of the tree fragments previously defined by marking the input tree, thereby obtaining the split tree. Therefore, the tree fragment is specified by one of the marks used in the split tree. The procedure terminates once a correctly marked node is found.

Procedure 50 commences with step 41 which accepts two input parameters:

- startNode - The node defining the subtree wherein the marked node is searched for; and
- mark - The mark that marks the node being searched for.

Procedure 50 proceeds to step 43 where it is determined whether 'startNode' is  
5 marked with 'mark'. If so, this is the node that is searched for, and the procedure 50  
terminates in step 44 by returning 'startNode' to its calling method or procedure.

If the 'startNode' is not marked with 'mark', step 46 follows which determines  
that startNode is marked with another mark, such indicates that 'startNode' and all the  
other nodes in its subtree are marked with a mark other than 'mark'. It is not possible for  
10 any node marked with 'mark' to be found in the subtree under review, and the procedure  
50 terminates in step 47 by returning 'NULL' to its calling method or procedure.

Therefore, only an unmarked node proceeds to step 48 from where, if  
'startNode' has only one child, procedure 50 is called in step 49 with the child node as  
'startNode'.

15 If the procedure 50 with the child node returns with a found node, the  
'startNode' belongs to the tree fragment marked with 'mark'. The procedure 50 returns  
'startNode' and terminate in step 52. If the procedure 50 returns with 'NULL', neither  
'startNode' nor any of the nodes in the subtree belong to the tree fragment marked with  
'mark'. Procedure 50 terminates after returning 'NULL' in step 53.

20 Therefore, only unmarked nodes with two children proceed to step 54 from  
where procedure 50 is called with the first child as 'startNode'. It is followed by step 55  
from where procedure 50 is called with the second child as 'startNode'.

If the procedure 50 returns with 'NULL' for both the first child, determined in  
step 56, and the second child, determined in step 57, the calling procedure 50 terminates  
25 after also returning 'NULL' in step 58. Thus 'mark' has not been found in any one of the

children. If 'mark' has been found in one of the children, procedure 50 will terminate after returning that child node to the calling method or procedure in step 60 or step 62. The remaining alternative is that both children were marked with 'mark' in which case the criteria in both step 56 and step 61 are not met. The procedure 50 terminates by  
5 returning 'startNode', being the parent of the two children, at step 63.

Referring to Fig. 12, procedure 70 for generating nodes of a tree fragment is presented. Procedure 70 is invoked recursively to find all the nodes of a tree fragment in a split tree and to generate a whole tree fragment. The tree fragment is defined by a specific mark.

10 In step 75, the procedure accepts three parameters:

- node - A node of a tree fragment;
  - mark - The mark that marks the nodes belonging to a tree fragment in a split tree;
- and
- perform - A function that is invoked for each node belonging to the tree fragment

15 defined by 'mark'.

If step 76 determines that 'node' is equal to NULL, there are no nodes in that subtree marked with 'mark' and step 77 returns the procedure to its calling method or procedure.

20 However, if 'node' is not equal to NULL in step 76 then step 80 invokes the 'perform' function. This caller-defined procedure performs required operations on the nodes of the 'mark' tree fragment. This may include saving the result to a file, modifying the attributes of the nodes or displaying the result in a tree form on a user interface.

If step 82 determines that 'node' has no children (a leave node), the procedure 70 terminates in step 77 and returns to its calling procedure or method. If step 84 determines  
25 that 'node' has only one child, the child subtree is searched in step 86 to find the next

node belonging to the tree fragment presently being generated. After finding the next node marked with 'mark', procedure 70 is called again in step 88 with this found next node as 'node'. Procedure 70 calling itself recursively will continue adding new child nodes to the generated tree until there are no more nodes in the subtree defined by the  
5 child node marked with mark.

If 'node' has two children the procedure 70 will proceed to step 90 there, the children subtrees are treated in the same way as a single child in that steps 90 and 91 will find all the nodes in the subtree defined by the first child and generate a tree until there are no more nodes in the subtree defined by the first child node marked with mark. Steps  
10 94 and 95 repeat identical steps for the second child. The procedure 70 terminates in step 77.

In the preferred embodiment of the invention, method 40, illustrated in Fig. 10, starts in step 100 by accepting three parameters:

- split\_tree - the root node of the split tree;
  - 15 • mark - the mark that marks the nodes belonging to a tree fragment in a split tree;
- and
- perform - The function that is invoked for each node belonging to the 'mark' tree fragment.

Method 40 proceeds to step 103 where a root node of a tree fragment in a split  
20 tree is found. This is performed by calling procedure 50, which will return the first node found with 'mark' as its marking. Procedure 50 may return 'NULL' if the split tree does not contain a tree fragment marked with 'mark'.

Having established the root node in step 103, step 105 calls procedure 70 passing the root node as 'node'. As discussed above, procedure 70 is a recursive procedure for

iterating over all the nodes of the tree fragment in the split tree, finding the correctly marked nodes and creating the required tree fragment. Method 40 terminates in step 106.

The foregoing preferred methods describe only a number of embodiments of the present invention, and modifications, can be made thereto without departing from the scope of the present invention as defined in the appended claims.

The method 30 of marking an input tree to create a split tree and method 40 of creating tree fragments from the split tree, are preferably practiced using a conventional general-purpose computer system 200, such as that shown in Fig. 13 wherein the processes of Figs. 9 to 12 may be implemented as software, such as an application program executing within the computer system 200. In particular, the steps of methods 30 and 40 are effected by instructions in the software that are carried out by the computer. The software may be stored in a computer readable medium, including the storage devices described below, for example. The software is loaded into the computer from the computer readable medium, and then executed by the computer. A computer readable medium having such software or computer program recorded on it is a computer program product. The use of the computer program product in the computer preferably effects an advantageous apparatus for 30 for marking an input tree to create a split tree and for creating tree fragments from the split tree in accordance with the embodiments of the invention.

The computer system 200 comprises a computer module 202, input devices such as a keyboard 210 and mouse 212, output devices including a printer 208 and a display device 104.

The computer module 202 typically includes at least one processor unit 214, a memory unit 218, for example formed from semiconductor random access memory (RAM) and read only memory (ROM), input/output (I/O) interfaces including a video



interface 217, and an I/O interface 216 for the keyboard 210 and mouse 212. A storage device 224 is provided and typically includes a hard disk drive 226 and a floppy disk drive 228. A magnetic tape drive (not illustrated) may also be used. A CD-ROM drive 220 is typically provided as a non-volatile source of data. The components 214  
5 to 228 of the computer module 202, typically communicate via an interconnected bus 230 and in a manner which results in a conventional mode of operation of the computer system 200 known to those in the relevant art. Examples of computers on which the embodiments can be practised include IBM-PC's and compatibles, Sun Sparcstations or alike computer systems evolved therefrom.

10 Typically, the application program of the preferred embodiment is resident on the hard disk drive 226 and read and controlled in its execution by the processor 214. Intermediate storage of the program may be accomplished using the semiconductor memory 218, possibly in concert with the hard disk drive 226. In some instances, the application program may be supplied to the user encoded on a CD-ROM or floppy disk  
15 and read via the corresponding drive 220 or 228. Still further, the software can also be loaded into the computer system 200 from other computer readable medium including magnetic tape, a ROM or integrated circuit, a magneto-optical disk, a radio or infra-red transmission channel between the computer module 202 and another device, a computer readable card such as a PCMCIA card, and the Internet and Intranets including email  
20 transmissions and information recorded on websites and the like. The foregoing is merely exemplary of relevant computer readable mediums. Other computer readable mediums may be practiced without departing from the scope and spirit of the invention.

The methods of marking an input tree to create a split tree and of creating tree fragments from the split tree may alternatively be implemented in dedicated hardware

such as one or more integrated circuits performing the functions or sub functions of processes 30 and/or 40.

### **Industrial Applicability**

It is apparent from the above that the embodiments of the invention are  
5 applicable to computer and data processing industries.

The foregoing describes only some embodiments of the present invention, and modifications and/or changes can be made thereto without departing from the scope and spirit of the as be illustrative and not restrictive.

In the context of this specification, the word "comprising" means "including  
10 principally but not necessarily solely" or "having" or "including" and not "consisting only of". Variations of the word comprising, such as "comprise" and "comprises" have corresponding meanings.

**Claims:**

1. A method of creating a split tree for representing an input tree and at least one tree fragment, wherein said input tree comprises a plurality of nodes, said  
5 method comprising the steps of:

determining which of said plurality of nodes fit into a galley target; and

marking said nodes that fit into said galley target with a mark to create one said tree fragment for each galley target, said marked input tree being said split tree.

10 2. A method of splitting a split tree into at least one tree fragments, wherein each one of said tree fragments is associated with a mark and said split tree comprises at least one node marked said mark, said method comprising the steps of:

identifying said at least one node with said mark; and

15 creating said at least one tree fragment from said nodes marked with each of said marks.

3. Apparatus for creating a split tree for representing an input tree and at least one tree fragment, wherein said input tree comprises a plurality of nodes, said apparatus comprises:

20 means for determining which of said plurality of nodes fit into a galley target; and

means for marking said nodes that fit into said galley target with a mark to create said tree fragment for each galley target, said marked input tree being said split tree.

25

4. Apparatus for splitting a split tree into at least one tree fragments, wherein each one of said tree fragments is associated with a mark and said split tree comprises at least one node marked said mark, said apparatus comprises:

means for identifying said at least one node with said mark; and

5 means for creating said at least one tree fragment from said nodes marked with each of said marks.

5. A computer program product including a computer readable medium incorporating a computer program for creating a split tree for representing an input tree  
10 and at least one tree fragment, wherein said input tree comprises a plurality of nodes, said computer program product comprising:

means for determining which of said plurality of nodes fit into a galley target;

and

means for marking said nodes that fit into said galley target with a mark to  
15 create said tree fragment for each galley target, said marked input tree being said split tree.

6. A computer program product including a computer readable medium incorporating a computer program for splitting a split tree into at least one tree  
20 fragments, wherein each one of said tree fragments is associated with a mark and said split tree comprises at least one node marked said mark, said computer program product comprising:

means for identifying said at least one node with said mark; and

means for creating said at least one tree fragment from said nodes marked with  
25 each of said marks.

7. A method of creating a split tree for representing an input tree and at least one tree fragment, said method substantially as described herein with reference to Figs 7 to 13.

5

DATED this twenty-first Day of June, 1999

Canon Kabushiki Kaisha  
Patent Attorneys for the Applicant

SPRUSON & FERGUSON

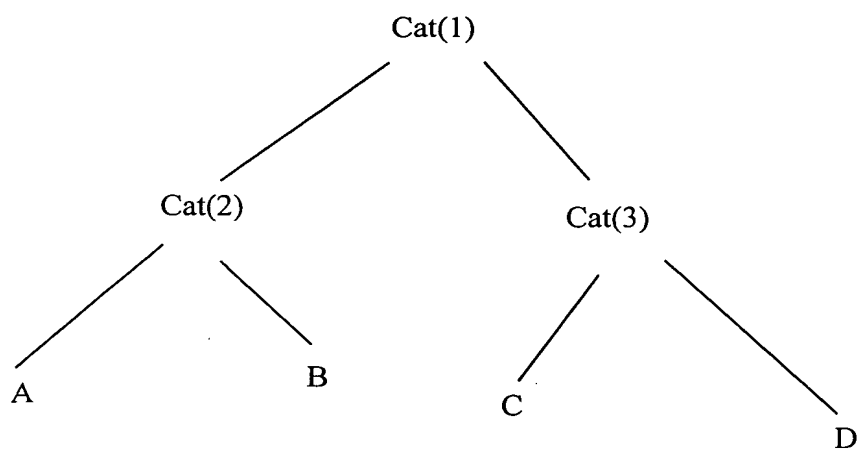


Fig. 1

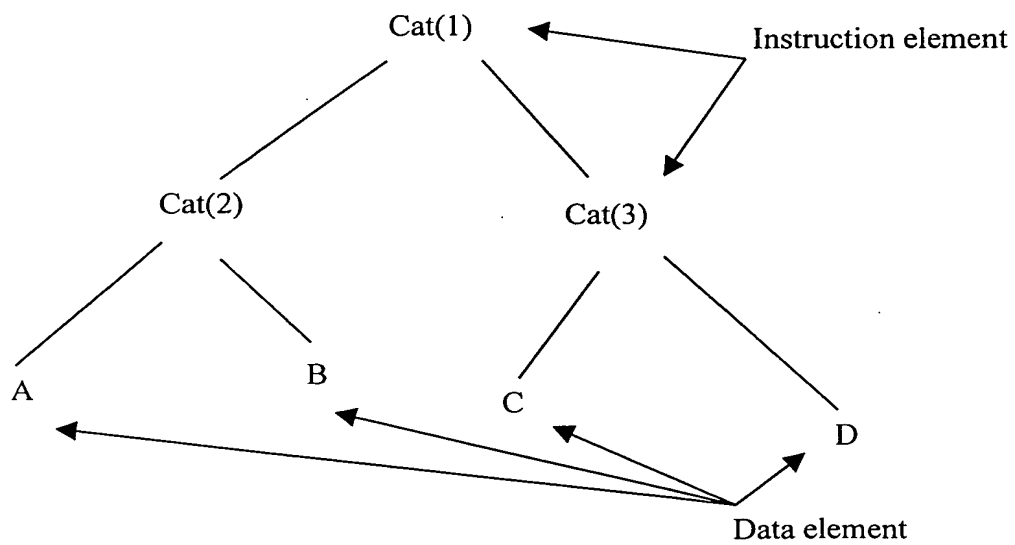


Fig. 2

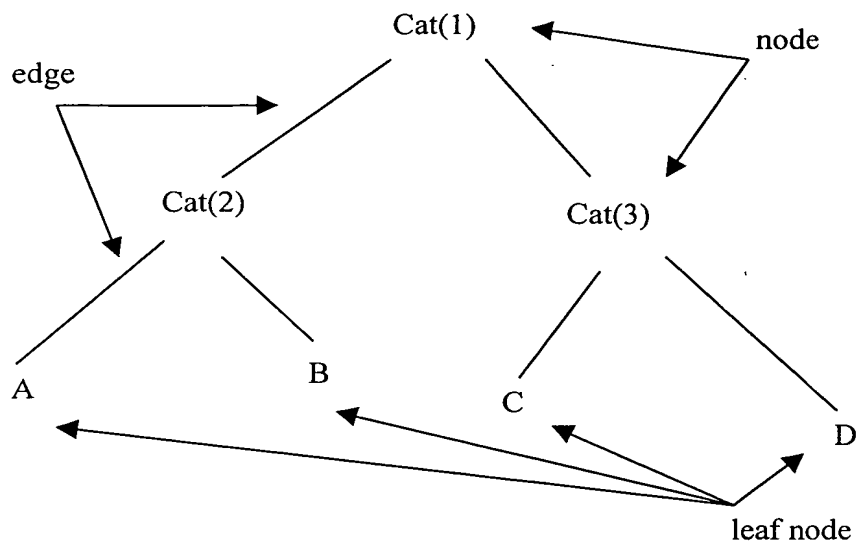


Fig. 3

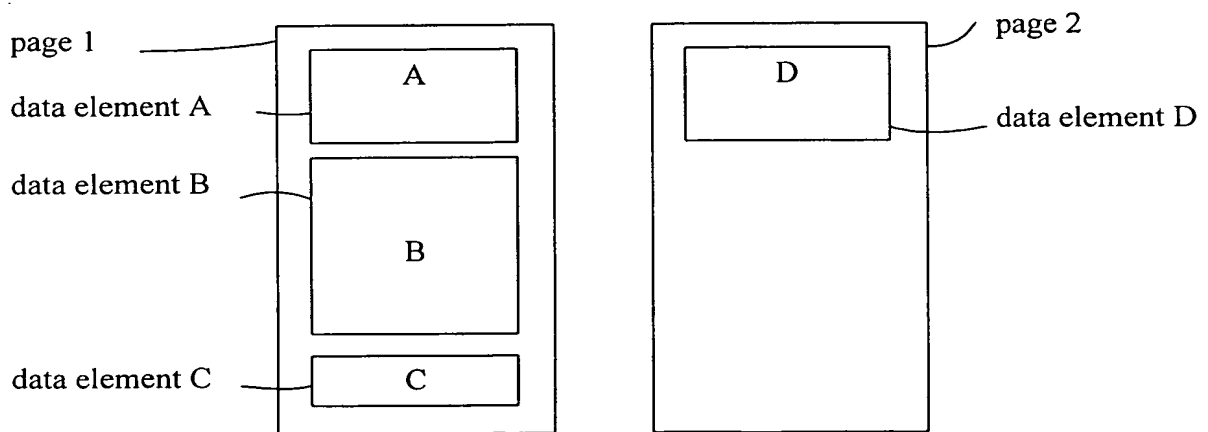


Fig. 4

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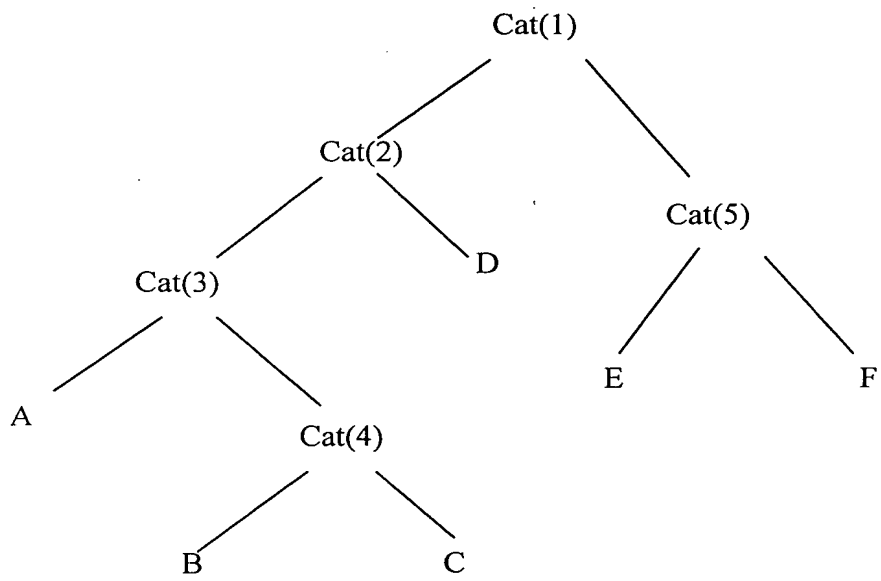


Fig. 5

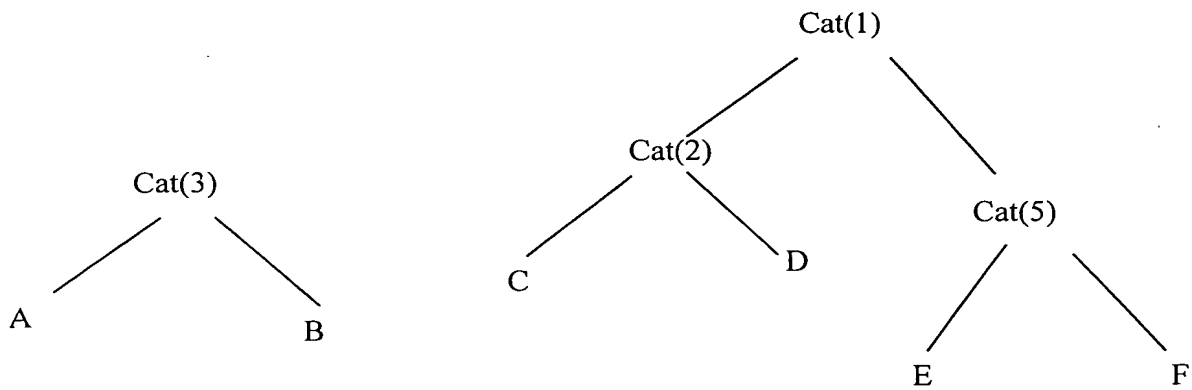


Fig. 6A

Fig. 6B



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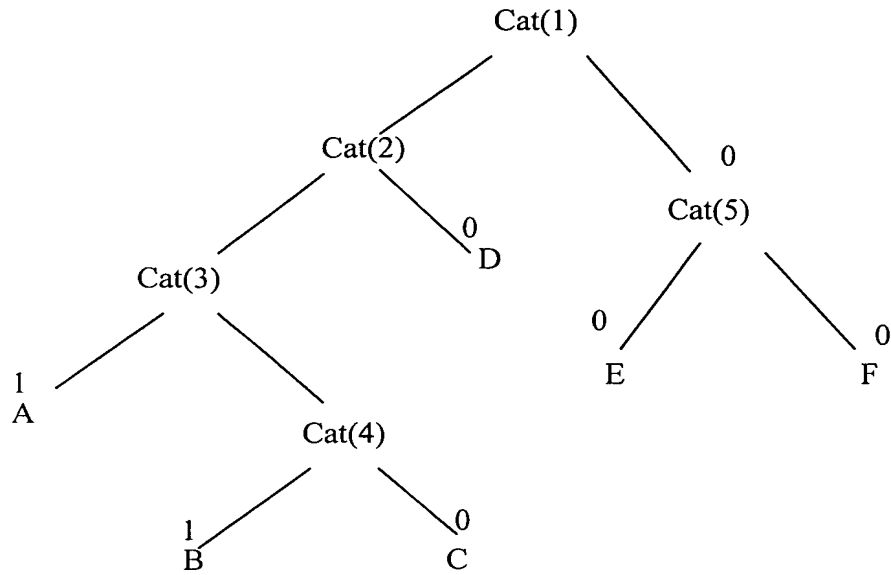


Fig. 7

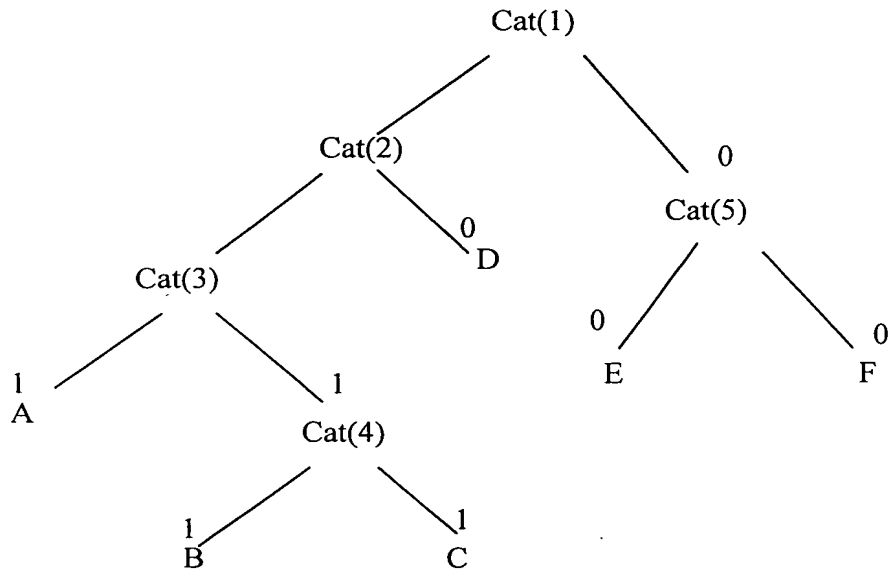


Fig. 8

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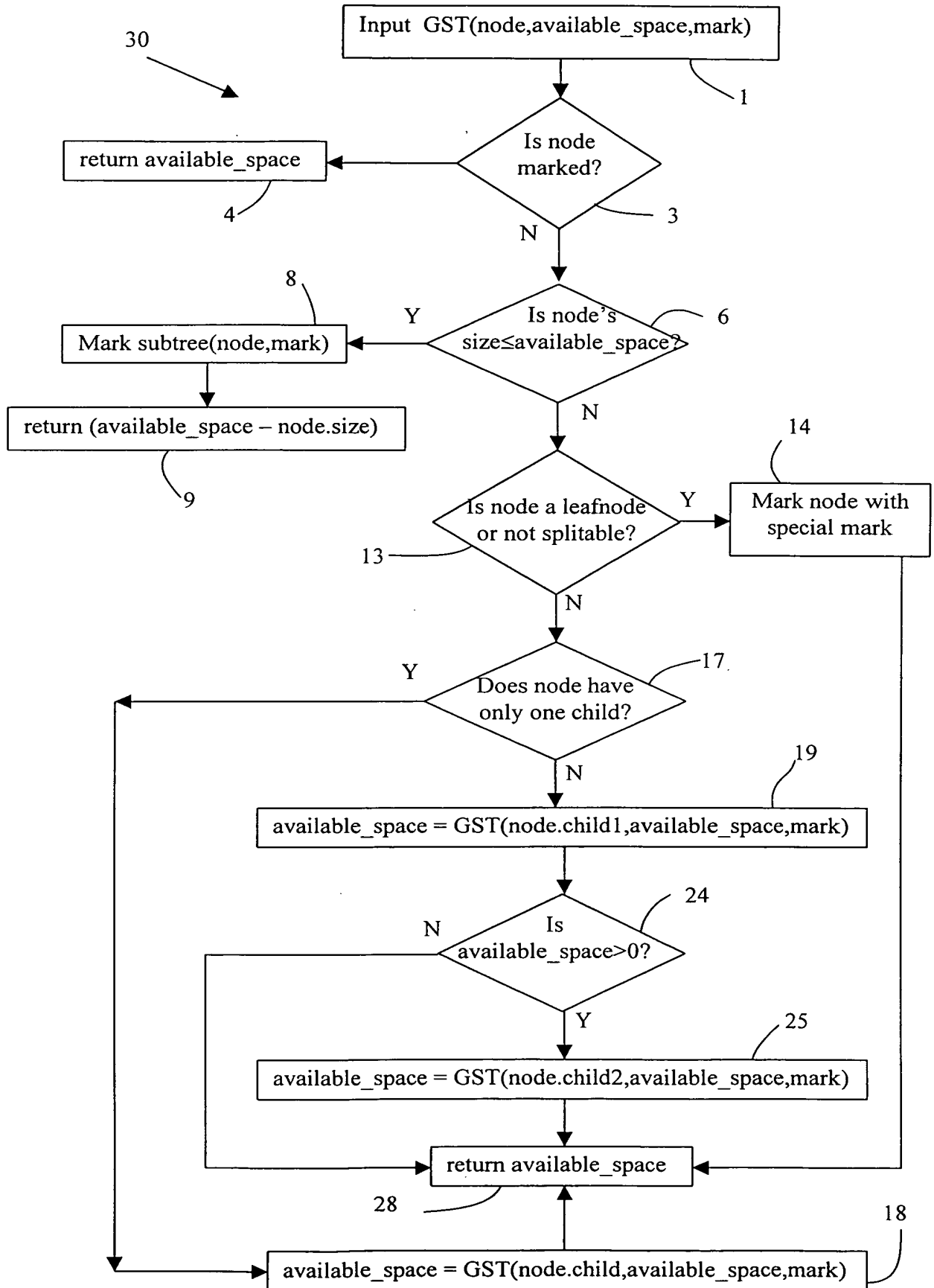


Fig. 9

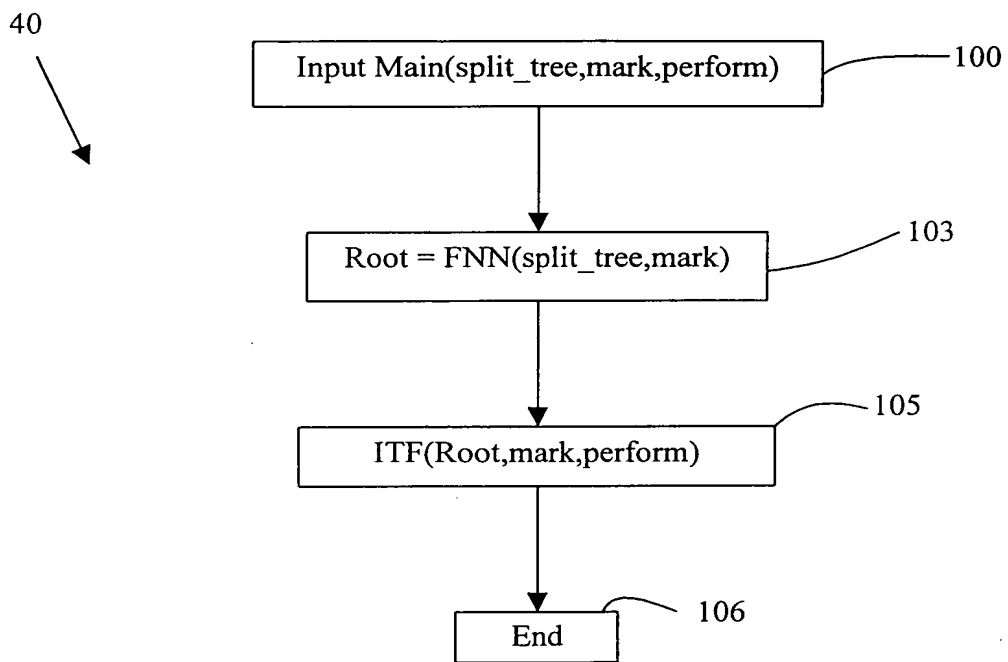


Fig. 10

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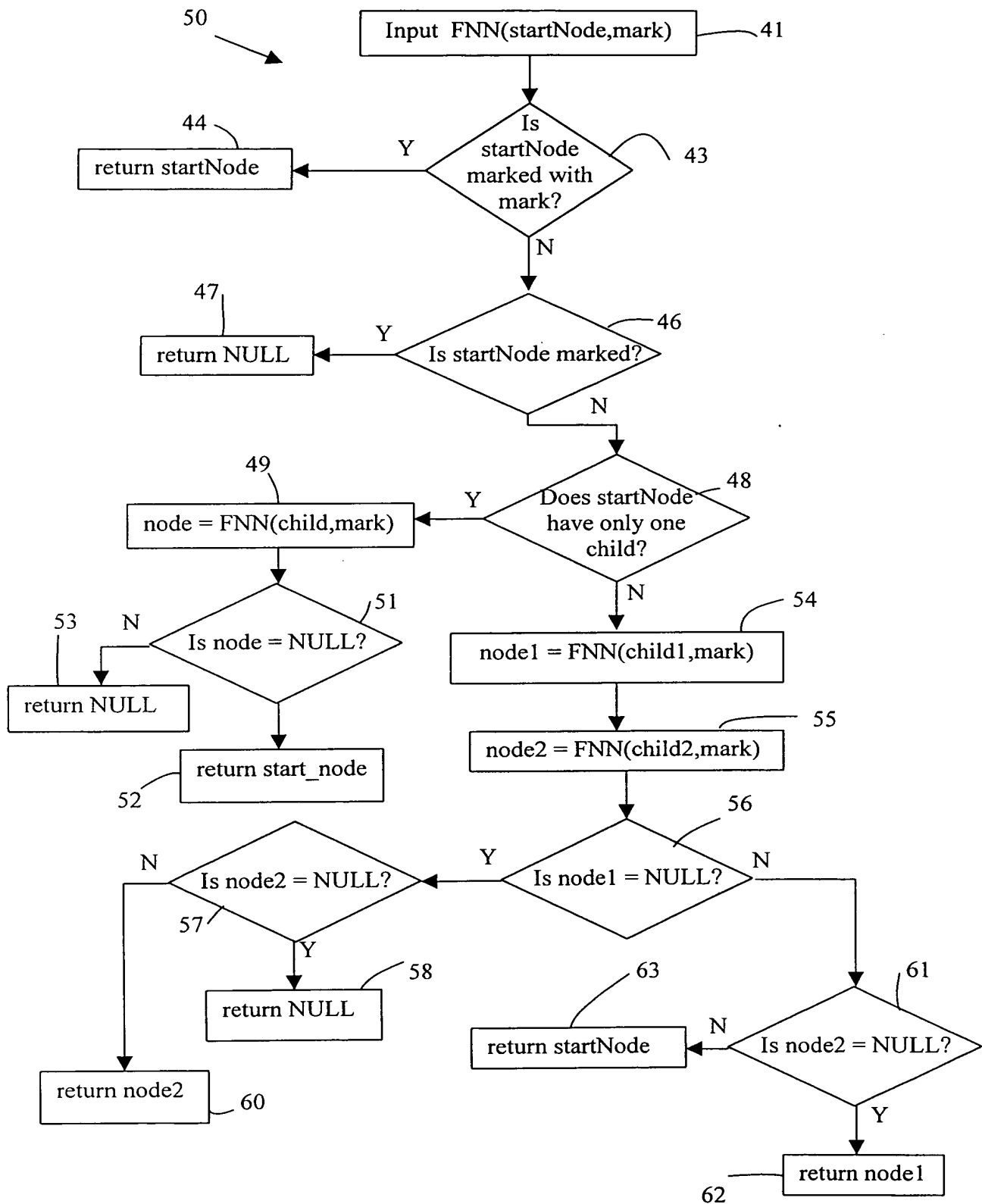


Fig. 11

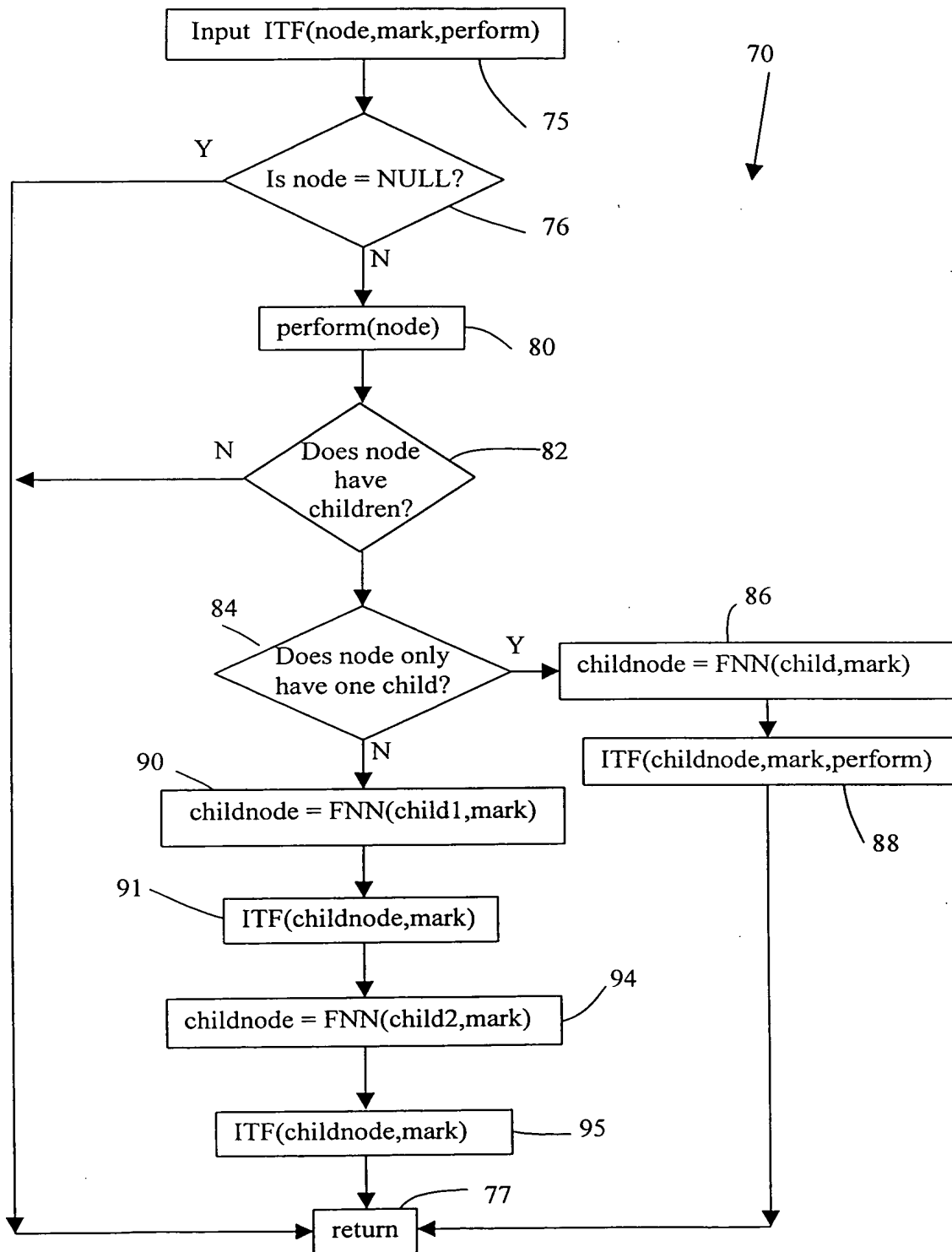


Fig. 12

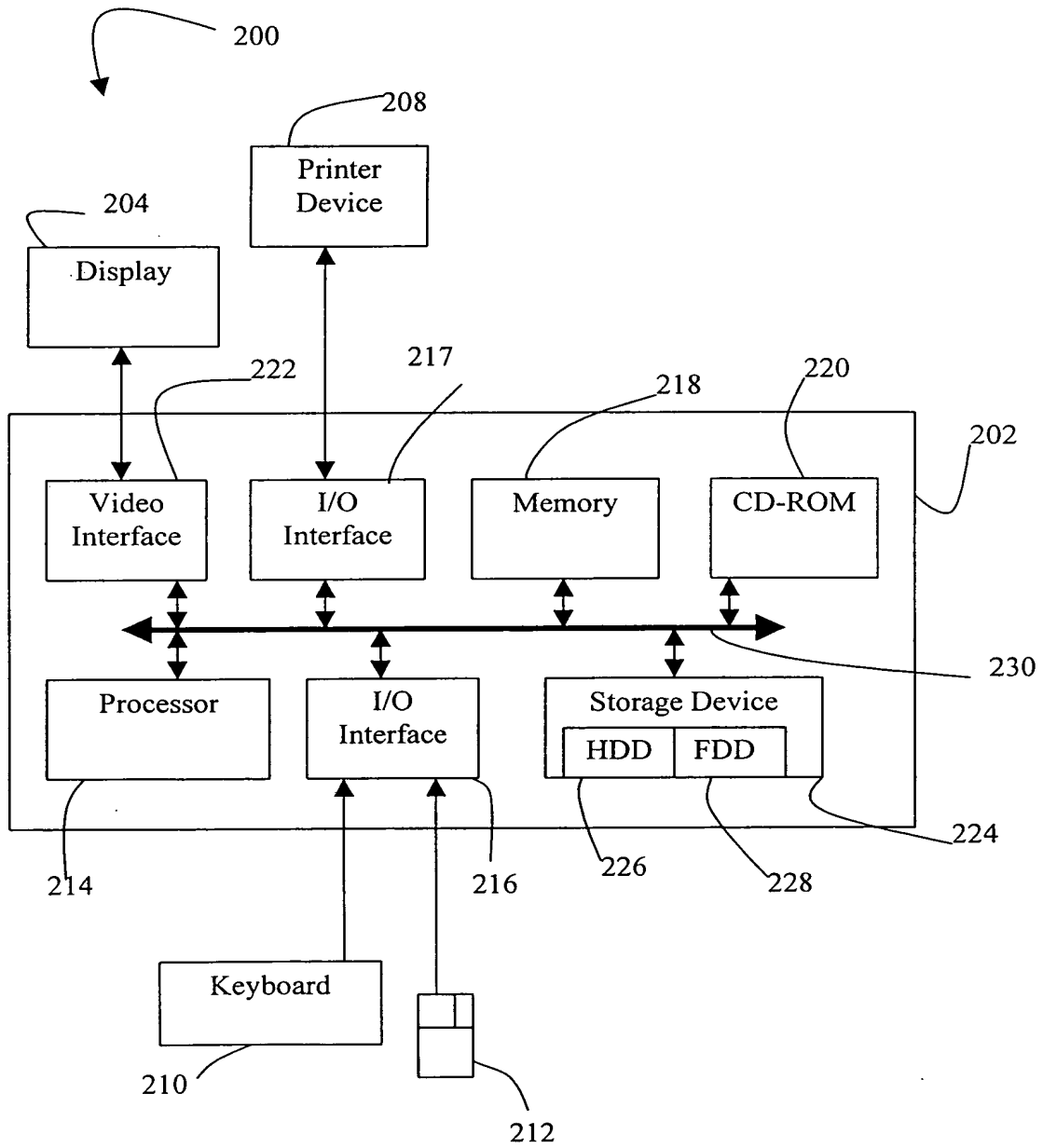


FIG. 13